

Nonbinary Hamming codes

$$n = \frac{q^r - 1}{q - 1}, \quad k = \frac{q^r - 1}{q - 1} - r, \quad d = 3$$

q -prime

Take $P(x) = a_0 + a_1 x + a_2 x^2 + \dots + a_{r-1} x^{r-1} + x^r$ - primitive

$$a_i \in \{0, 1, \dots, q-1\} = \mathbb{Z}_q$$

Construct $\mathbb{Z}_q^r = \{0, 1, \alpha, \alpha^2, \dots, \alpha^{q^r-2}\}$

where $P(\alpha) = 0 \quad \alpha = 00\dots 010$.

* Take $K = [1 \ \alpha \ \alpha^2 \ \alpha^3 \ \dots \ \alpha^{n-1}]$

$$n = \frac{q^r - 1}{q - 1}$$

all columns
are different
and $\alpha^i \neq a \alpha^j$
 $a \in \mathbb{Z}_q$

PROOF:
 Let $\alpha^i = a \alpha^j$ $a \in \mathbb{Z}_q$
 $0 \leq i, j \leq n-1$, $n = \frac{q^r - 1}{q - 1}$

let $j > i$
 Then $\alpha^{j-i} = a$

let $s = j - i$ then $s < n$
 $\alpha^s = a$ (α -primitive)

$\alpha^{s(q-1)} = a^{q-1} = 1$ but by Fermat Theorem

$s(q-1) < n \cdot (q-1) = q^r - 1$

Contradiction

$$\begin{cases} \alpha^{s(q-1)} = 1 \\ s(q-1) < q^r - 1 \end{cases}$$

Example $q=3$ $r=2$

$$n = \frac{q^2 - 1}{q - 1} = q + 1 = 4. \quad d=3$$

$$P(x) = x^2 + x + 2$$

$$\mathbb{Z}_3^2 = \mathbb{Z}_9^r$$

$$\mathbb{Z}_9 = \{0, 1, 2\}$$

$$\alpha^2 + \alpha + 2 = 0$$

$$\alpha^2 = 2\alpha + 1$$

0	0	0
0	1	1
0	2	α^4
1	0	α
1	1	α^7
1	2	α^6
2	0	α^5
2	1	α^2
2	2	α^3
α	1	$\alpha^8 = 1$

$$H = [1 \ \alpha \ \alpha^2 \ \alpha^3] =$$

$$= \begin{bmatrix} 0 & 1 & 2 & 2 \\ 1 & 0 & 1 & 2 \end{bmatrix}$$

Consider $v = (2110)$

$$\text{Then } Hv = [1 \ \alpha \ \alpha^2 \ \alpha^3] \begin{bmatrix} 2 \\ 1 \\ 1 \\ 0 \end{bmatrix} =$$

$$= 2 + \alpha + \alpha^2 = 0 \iff P(\alpha) = 0$$

(2110) is a codeword.

Let $v = (v_1 \ v_2 \ v_3 \ v_4)$

$$\text{then } v(x) = v_1 + x v_2 + x^2 v_3 + x^3 v_4$$

$$Hv = v_1 + \alpha v_2 + \alpha^2 v_3 + \alpha^3 v_4 = 0 \Rightarrow$$

v is a codeword iff

$$v(\alpha) = 0 \quad \alpha \text{ is a root of } v(x)$$

If $v \in \mathbb{C} \Rightarrow v(\alpha) = 0$

$w(x) = v(x) Q(x)$ For any $Q(x)$

$w(\alpha) = v(\alpha) Q(\alpha) = 0 \Rightarrow w \in \mathbb{C}$

code consists of all multiples
of $v(x)$ if $v \in \mathbb{C}$

$$x^n = 1$$

Multiplication by x is equivalent
to rotation (cyclic shift)

example $v = (2110)$

$$v(x) = 2 + x + x^2$$

$$w(x) = \text{Rot } v(x) = 2x + x^2 + x^3$$

$$w = 0211$$

$$y = \text{Rot } w = 1021$$

$$y(x) = 1 + 2x^2 + x^3 = w(x) \cdot x =$$

$$\begin{aligned}
 &= (2x + x^2 + x^3) x = 2x^2 + x^3 + x^4 = \\
 &= 1 + 2x^2 + x^3 \Rightarrow \quad \quad \quad \uparrow \\
 & \quad \quad \quad \quad \quad \quad \quad \quad \quad \quad \quad x^4 = 1
 \end{aligned}$$

$$y = 1021.$$

Since $P(\alpha) = 0 \quad P \in \mathbb{C}$

For our example

$$P(x) = x^2 + x + 2 \Rightarrow$$

$$P = (2 \ 1 \ 1 \ 0) \in \mathbb{C}$$

$$w = \text{Rot } P = (0 \ 2 \ 1 \ 1) \in \mathbb{C}$$

$\text{Rot } w = (1 \ 0 \ 2 \ 1) \in \mathbb{C}$ to verify

$$[1 \ \alpha \ \alpha^2 \ \alpha^3] \begin{bmatrix} 1 \\ 0 \\ 2 \\ 1 \end{bmatrix} = 1 + 2\alpha^2 + \alpha^3$$

$$\alpha^2 = 2\alpha + 1$$

$$\alpha^3 = 2\alpha + 2$$

Thus

$$1 + 2\alpha^2 + \alpha^3 =$$

$$= 1 + 2(2\alpha + 1) + 2\alpha + 2 =$$

$$= 1 + 6\alpha + 5 = 0 \quad \text{Thus}$$

$$1021 = \text{Rot}(0211) \in \mathbb{C}$$

Thus we constructed

$(4, 3, 3)$ ternary cyclic
Hamming code with

the check matrix

$$H = [1 \alpha \alpha^2 \alpha^3] = \begin{bmatrix} 0 & 1 & 2 & 2 \\ 1 & 0 & 1 & 2 \end{bmatrix} r=2$$

the generating matrix for this code is

$$G = P(x) = [2 \ 1 \ 1 \ 0]$$

EXAMPLE: $q=3 \ r=3$

$$n = \frac{q^3 - 1}{q - 1} = 13$$

$$k = 13 - 3 = 10$$

$$H = [1 \ \alpha \ \alpha^2 \ \alpha^3 \ \alpha^4 \ \alpha^5 \ \alpha^6 \ \alpha^7 \ \alpha^8 \ \alpha^9 \ \alpha^{10} \ \alpha^{11} \ \alpha^{12}]$$

$$G = \begin{bmatrix} P(x) \\ xP(x) \\ x^2P(x) \\ x^3P(x) \\ x^4P(x) \\ x^5P(x) \\ x^6P(x) \\ x^7P(x) \\ x^8P(x) \end{bmatrix}$$